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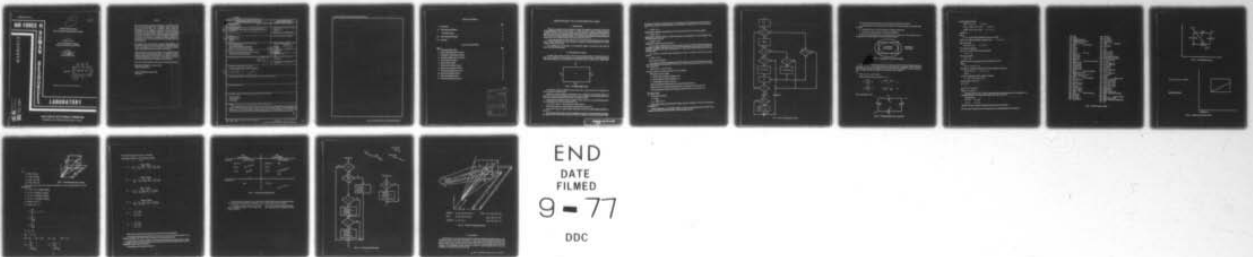
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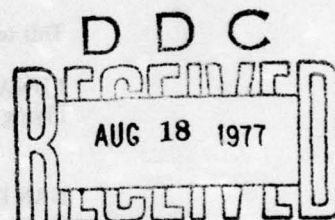
**COMPUTER GRAPHICS:
TWO- AND THREE-DIMENSIONAL CLIPPING**

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May 1977



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COMPUTER GRAPHICS: TWO- AND THREE-DIMENSIONAL CLIPPING

I. INTRODUCTION

Clipping is a procedure used to determine the proper line segment(s) of a given line, such that the resulting segment(s) contain no points exterior to a plane that serves to define a boundary. Clipping is employed to limit the amount of work a display device must perform by drawing only the visible parts of a line. This is done by using a clipping algorithm to compute only the line segments interior (or exterior) to a defined region that has been selected for display. The resulting list of lines represents a substantially smaller amount of work for the display hardware to process in a given period of time.

A good clipping algorithm quickly rejects lines that lie outside the viewing area, and since speed is always desirable, algorithms that lend themselves to parallel (or pipeline implementation), in either hardware or firmware, are sought.

Let us initially limit our discussion to two-dimensional clipping. A later section of this paper will explore three-dimensional clipping.

II. TWO-DIMENSIONAL CLIPPING

Typically, display regions have been chosen as being rectangular (Figure 1). The actual number of sides is not of importance to the development of a clipping algorithm, although it does affect performance when finally implemented. Viewing windows are always chosen however, as convex sets in a Hilbert Space.

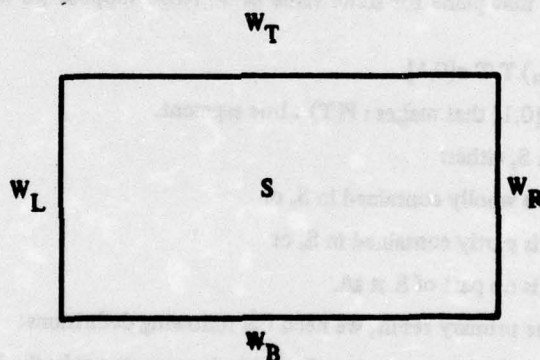


Figure 1. Rectangular display region.

The standard properties of Hilbert spaces, as well as the set concepts referred to in this paper, are found in most texts on linear analysis.

Making use of these properties, the following is clear: A straight line, intersecting the interior of a convex set, can intersect the boundary of the set in, at most, two places. In the event the convex set is closed and bounded, the straight line will intersect it in exactly two places.

At this point, we introduce the idea of "inward" normals. That is, we say n is an inward normal for any boundary point X of S if, for any other point Y in S ,

$$n \cdot (Y-X) > 0, \text{ where } "\cdot" \text{ denotes the inner product operation.}$$

With this representation, we have defined a natural geometry over the set S .

For the set S only four unique inward normals exist. All others are mathematically equivalent to these four.

Our particular clipping problem, that of determining the points of intersection between a particular line segment and the set S can then be partially resolved by examining the parameterized behavior of the

line segment with the four inward normals to S . In particular, if f is a boundary point of S for which n is an inward vector and $P(T)$ is a parameterization of the line segment in question, then, for particular values of T ,

$$n \cdot [P(T) - f] > 0$$

implies $P(T)$ is "pointed" toward the interior of S with respect to at least one of the four normals,

$$n \cdot [P(T) - f] = 0$$

implies $P(T)$ is "pointed" parallel to the plane containing f and perpendicular to n (that is, parallel to this particular boundary of S) and

$$n \cdot [P(T) - f] < 0$$

implies $P(T)$ is "pointed" away from the set S .

So far we have mentioned only that any line piercing S does so at exactly two points. It is also clear that these two points of intersection do not lie on the same target planes to the boundary of S (a plane not completely containing a line may be intersected by it at most once). These facts give rise to our first important result:

if f is a point in a particular plane bounding S for which n is an inward normal, then the scalar equation

$$n \cdot [P(T) - f] = 0$$

has, at most, one solution.

Thus, there are exactly two possibilities; first, $P(T)$ is parallel to some plane lying along the boundary of S . Second $P(T)$ pierces that plane for some value of T . Now, suppose we want to talk about drawing only a line segment:

$$P(T) = P_0 + (P_1 - P_0)T \quad T \in [0,1]$$

It is the restriction of T to $[0,1]$ that makes $P(T)$ a line segment.

With regard to the set S , either:

1. The line segment is wholly contained in S , or
2. The line segment is partly contained in S , or
3. The line segment is no part of S at all.

Before introducing the primary result, we need the following definitions:

Define S_b as the set of all planes bounding S , that is, intersecting only the boundary of S .

For each $b \in S_b$, define $n(b)$ as the set of positive normals drawn from b with respect to S .

The Primary Result

The set of all T such that:

1. $T \in (0,1)$
2. $n \cdot [P(T) - f] > 0$

Where for every $b \in S_b$, $n \in n(b)$, and f is an arbitrary vector in b , completely determines the line segment containment in S .

Careful examination of the flowchart (Figure 2) reveals the underlying geometry involved in this technique. The implementation lends itself to swift line segment rejection.

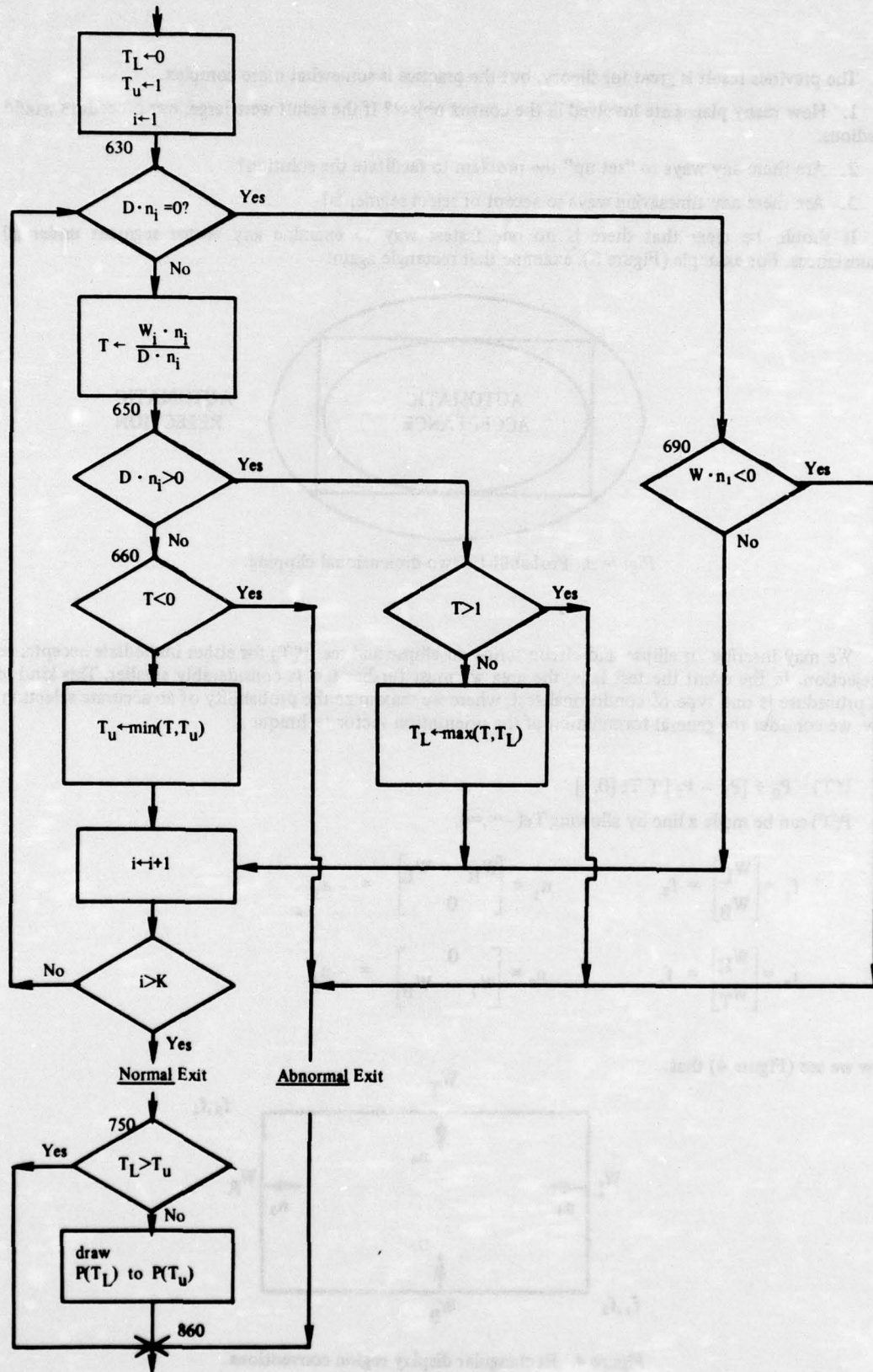


Figure 2. Primary climbing algorithm flowchart.

The previous result is great for theory, but the practice is somewhat more complex.

1. How many planes are involved in the convex object? If the result were large, our procedure would be tedious.

2. Are there any ways to "set up" the problem to facilitate the solution?

3. Are there any timesaving ways to accept or reject segments?

It should be clear that there is no one fastest way to examine any vector segment under all circumstances. For example (Figure 3), examine that rectangle again:

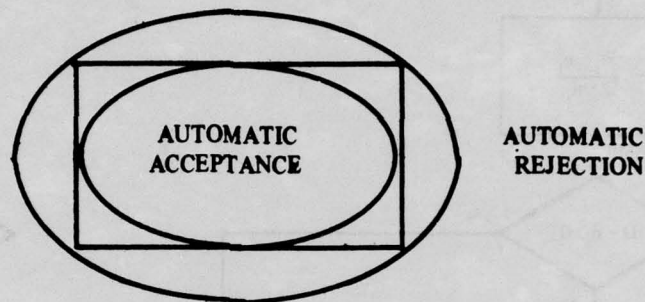


Figure 3. Probabilistic two-dimensional clipping.

We inscribe an ellipse and circumscribe an ellipse and test $P(T)$ for either immediate acceptance or rejection. In the event the test fails, the area we must further test is considerably smaller. This kind of test procedure is one type of conditional test, where we maximize the probability of an accurate selection. Now, we consider the general formulation of the orientation vector technique.

Let

$$P(T) = P_0 + [P_1 - P_0] T \quad T \in [0,1]$$

$P(T)$ can be made a line by allowing $T \in (-\infty, \infty)$.

$$\begin{aligned} f_1 &= \begin{bmatrix} W_L \\ W_B \end{bmatrix} = f_2 & n_1 &= \begin{bmatrix} W_R - W_L \\ 0 \end{bmatrix} = -n_3 \\ f_3 &= \begin{bmatrix} W_R \\ W_T \end{bmatrix} = f_4 & n_2 &= \begin{bmatrix} 0 \\ W_T - W_B \end{bmatrix} = -n_4 \end{aligned}$$

Now we see (Figure 4) that:

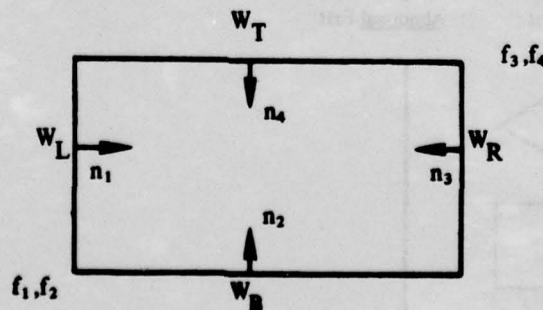


Figure 4. Rectangular display region conventions.

The five conditions become

$$[P(T) - f_i] \cdot n_i > 0 \quad i = 1, 2, 3, 4$$

$$\text{or } P_0 \cdot n_i + T[P_1 - P_0] \cdot n_i > 0 \quad i = 1, 2, 3, 4$$

$$\text{or } T[P_1 - P_0] \cdot n_i > [f_i - P_0] \cdot n_i \quad i = 1, 2, 3, 4$$

and $T \in [0, 1]$

letting

$$P_1 - P_0 = D \text{ (D is called the Directorix of the vector } P(T))$$

$$\text{and } W_i = P_0 - f_i \text{ (} W_i \text{ is the } i\text{th plane "weight" vector), } i = 1, 2, 3, 4$$

Then we can write:

$$TD \cdot n_i + W_i \cdot n_i > 0 \quad i = 1, 2, 3, 4$$

Now we have our algorithm:

For a K-sided convex object S with inside normals:

$$n_1, n_2, \dots, n_K$$

Points on each plane:

$$f_1, f_2, \dots, f_K$$

Let $P(T)$ be the parameterization of the line running from P_0 to P_1

Define:

$$W_i = P_0 - f_i, D = P_1 - P_0$$

Then the K + 1 conditions for any portion of the line segment being in S are:

$$TD \cdot n_i + W_i \cdot n_i > 0 \quad i = 1, 2, \dots, K$$

$$T \in [0, 1]$$

Let the set of such that each ith case above is satisfied:

$$T_i = \{ T \mid TD \cdot n_i + W_i \cdot n_i > 0 \}$$

Then for each T_i , $T_i \in (-\infty, \infty)$

$$\text{Let } T = \bigcap_{i=1}^K T_i$$

$$i = 1$$

$$\text{Then } T = \{ T \mid P(T) \in S \}$$

$$\text{Finally, } T^* = T \cap [0, 1]$$

Using the Tektronix 4051 Graphic Computing System, an implementation of this algorithm was programmed (Figure 5). The program and sample results follows (Figure 6):

In this program:

$$P_{0x} = X1 \quad P_{1x} = X2$$

$$P_{0y} = Y1 \quad P_{1y} = Y2$$

Figures 7 and 8 show examples of segments drawn and rejected, respectively.

We can introduce efficiency into the computation by first evaluating directly for the values of T.


```

100 INIT
110 PAGE
120 WINDOW 0,22,0,22
130 VIEWPORT 65,130,0,65
140 DIM F(4,2),N(4,2)
150 REM ..... Define window
160 W1=7
170 W2=18
180 W3=10
190 W4=17
200 REM ..... Draw axis
210 MOVE 0,0
220 DRAW 22,0
230 MOVE 0,0
240 DRAW 0,22
250 REM ..... Draw Window boundary
260 MOVE W1,W3
270 DRAW W2,W3
280 DRAW W2,W4
290 DRAW W1,W4
300 DRAW W1,W3
310 REM ..... Enter line definition
320 PRINT @32,21,0,90
330 PRINT "Enter x1,y1, x2,y2 :";
340 INPUT X1,Y1,X2,Y2
350 REM ..... Normals defined
360 N(1,1)=W2-W1
370 N(1,2)=0
380 N(3,1)=-N(1,1)
390 N(3,2)=-N(1,2)
400 N(2,1)=0
410 N(2,2)=W4-W3
420 N(4,1)=-N(2,1)
430 N(4,2)=-N(2,2)
440 REM ..... Window boundary
450 F(1,1)=W1
460 F(1,2)=W3
470 F(2,1)=F(1,1)
480 F(2,2)=F(1,2)
490 F(3,1)=W2
500 F(3,2)=W4
510 F(4,1)=F(3,1)
520 F(4,2)=F(3,2)
530 REM ..... Algorithm
540 D1=X2-X1
550 D2=Y2-Y1
560 L=0
570 U=1
580 FOR I=1 TO 4
590 W1=X1-F(I,1)
600 W2=Y1-F(I,2)
610 D=D1*N(I,1)+D2*N(I,2)
620 W=W1*N(I,1)+W2*N(I,2)
630 IF D=0 THEN 690
640 T=-W/D
650 IF D>0 THEN 710
660 IF T<0 THEN 860
670 U=T MIN U
680 GO TO 730
690 IF W<0 THEN 680
700 GO TO 730
710 IF T>1 THEN 860
720 L=T MAX L
730 NEXT I
740 REM ..... Segment drawn
750 IF L=>U THEN 860
760 X4=X1+(X2-X1)*I
770 Y4=Y1+(Y2-Y1)*I
780 X5=X1+(X2-X1)*U
790 Y5=Y1+(Y2-Y1)*U
800 MOVE X4,Y4
810 DRAW X5,Y5
820 PRINT @32,21:0,30
830 PRINT "SEGMENT DRAWN"
840 END
850 REM ..... Segment rejected
860 PRINT @32,21:0,30
870 PRINT "SEGMENT REJECTED"
880 END

```

Figure 5. Sample computer program.

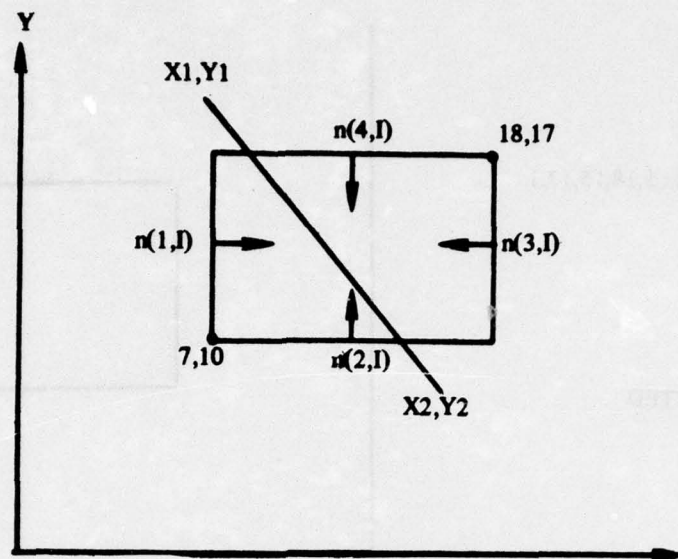
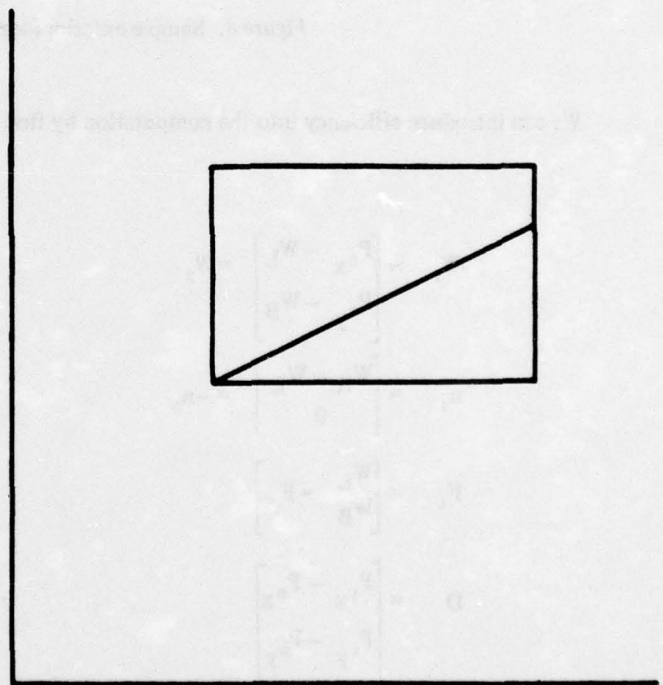


Figure 6. Sample display problem.

Enter x1,y1, x2,y2 : 7,10,20,16

SEGMENT DRAWN

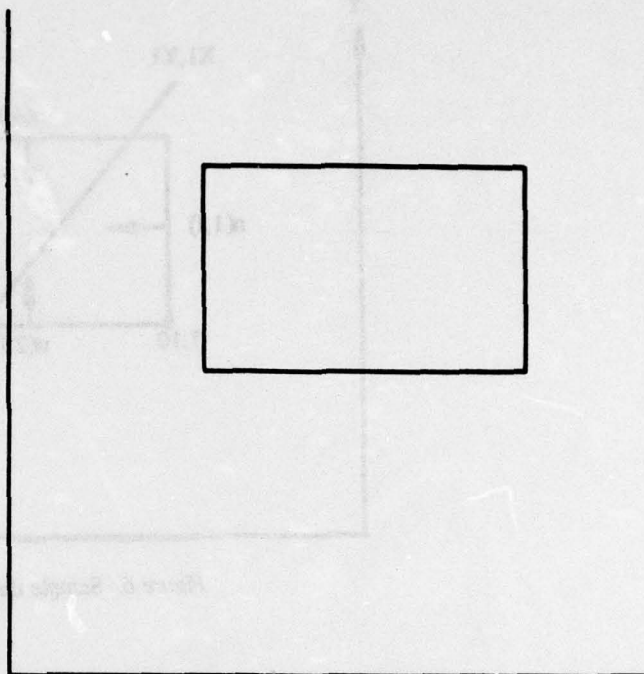


EXAMPLE #1

Figure 7. Sample interior (segment drawn)

Enter x1,y1, x2,y2 : 5,18,15,17.1

SEGMENT REJECTED



EXAMPLE #2

Figure 8. Sample exterior (segment rejected).

We can introduce efficiency into the computation by first evaluating directly for the values of T.

$$w_1 = \begin{bmatrix} P_{0x} - w_L \\ P_{0y} - w_B \end{bmatrix} = w_2$$

$$n_1 = \begin{bmatrix} w_R - w_L \\ 0 \end{bmatrix} = -n_3$$

$$F_1 = \begin{bmatrix} w_L \\ w_B \end{bmatrix} = F_2$$

$$D = \begin{bmatrix} P_{1x} - P_{0x} \\ P_{1y} - P_{0y} \end{bmatrix}$$

$$w_3 = \begin{bmatrix} P_{0x} - w_R \\ P_{0y} - w_T \end{bmatrix} = w_4$$

$$n_2 = \begin{bmatrix} 0 \\ w_T - w_B \end{bmatrix} = -n_4$$

$$F_3 = \begin{bmatrix} w_R \\ w_T \end{bmatrix} = F_4$$

- (1) $T(P_{1x} - P_{0x})(W_R - W_L) > (W_L - P_{0x})(W_R - W_L)$
- (2) $T(P_{1y} - P_{0y})(W_T - W_B) > (W_B - P_{0y})(W_T - W_B)$
- (3) $T(P_{1x} - P_{0x})(W_L - W_R) > (W_R - P_{0x})(W_L - W_R)$
- (4) $T(P_{1y} - P_{0y})(W_R - W_T) > (W_T - P_{0y})(W_B - W_T)$
- (5) $T \in [0,1]$

Assuming that $P(T)$ is not parallel to any side of the rectangle, then we can solve for the values of T so that each equation (1 through 4) is satisfied identically.

$$(1) = \frac{W_L - P_{0x}}{P_{1x} - P_{0x}}$$

$$(2) = \frac{W_B - P_{0y}}{P_{1y} - P_{0y}}$$

$$(3) = \frac{W_R - P_{0x}}{P_{1x} - P_{0y}}$$

$$(4) = \frac{W_T - P_{0y}}{P_{1y} - P_{1y}}$$

III. THREE-DIMENSIONAL CLIPPING

Since the algorithm was not restricted to two dimensions, the basic result applies in the three-dimensional case. We will examine a particular case. The set of points chosen for perspective purposes is centered with the axis center at the eye point, looking along the Z axis. The set of points is given by the set product

$$\{W_L, W_R\} \times \{W_B, W_T\} \times \{W_H, W_Y\}$$

The region in space R , now being viewed as our convex set, is given by the product

$$[W_L, W_R] \times [W_B, W_T] \times [W_H, W_Y]$$

In the 2-D case, there were four equations in T and one additional constraint. In this 3-D case, there are six equations in T and that same constraint.

There are six bounding planes, four of which contain the point $[0,0,0]$. The other two are the hither and yon planes, which are handled as special cases (Figure 9). Symmetry conditions imply the following:

$$W_L = -W_R \quad W_B = -W_T$$

Let

$$V_1 = [W_R, W_T, W_H]$$

$$V_2 = [-W_R, +W_T, W_H]$$

$$V_3 = [-W_R, -W_T, W_H]$$

$$V_4 = [+W_R, -W_T, W_H]$$

These four vectors are the vectors along the intersection of the four bounding planes to the first perspective plane.

$$n_1 = V_1 \times V_2 = [0, -2W_R W_H, 2W_R W_T]$$

$$n_2 = V_2 \times V_3 = [2W_T W_H, 0, 2W_R W_T]$$

$$n_3 = V_3 \times V_4 = [0, 2W_R W_H, 2W_R W_T]$$

$$n_4 = V_4 \times V_1 = [-2W_T W_H, 0, 2W_R W_T]$$

$$n_5 = (H_{\text{ither}}) = [0, 0, 1]$$

$$n_6 = (Y_{\text{on}}) = [0, 0, -1]$$

$$f_1 = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} = f_2 = f_3 = f_4$$

$$f_5 = \begin{bmatrix} 0 \\ 0 \\ W_H \end{bmatrix}$$

$$f_6 = \begin{bmatrix} 0 \\ 0 \\ W_Y \end{bmatrix}$$

$$D = P_1 - P_0$$

$$W_L = -f_1 + P_0$$

$$W_1 = +P_0$$

$$W_2 = +P_0$$

$$W_3 = +P_0$$

$$W_4 = +P_0$$

$$W_5 = \begin{bmatrix} +P_{0x} \\ +P_{0y} \\ -W_H + P_{0z} \end{bmatrix}$$

$$W_6 = \begin{bmatrix} +P_{0x} \\ +P_{0y} \\ -W_Y + P_{0z} \end{bmatrix}$$

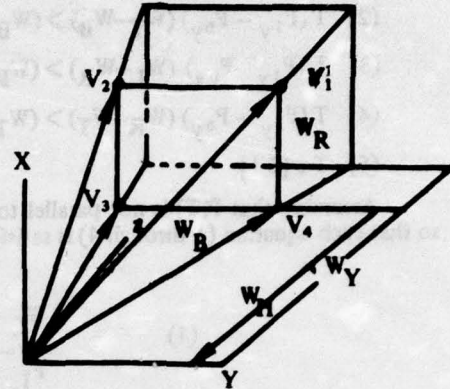


Figure 9. Three-dimensional viewing window.

The same rules apply to $D \cdot n_i$ and $w_i \cdot n_i$ as before.

The primary formula is $T = -\frac{w_i \cdot n_i}{D \cdot n_i}$ for each i , giving

$$T_1 = \frac{W_T P_{0z} + W_H P_{0y}}{(P_{1z} - P_{0z}) W_T - (P_{1y} - P_{0y}) W_H}$$

$$T_2 = \frac{W_H P_{0x} + W_R P_{0z}}{(P_{1x} - P_{0x}) W_H + (P_{1z} - P_{0z}) W_R}$$

$$T_3 = -\frac{W_H P_{0y} + W_T P_{0z}}{(P_{1y} - P_{0y}) W_H + (P_{1z} - P_{0z}) W_T}$$

$$T_4 = \frac{-W_R P_{0z} + W_H P_{0x}}{(P_{1z} - P_{0z}) W_R - (P_{1x} - P_{0x}) W_H}$$

$$T_5 = \frac{-P_{0z} + W_H}{P_{1z} - P_{0z}}$$

$$T_6 = \frac{-W_y + P_{0z}}{P_{0z} - P_{1z}}$$

These values of T , define the intersection points of $P(T)$ with each bounding plane.

It can be seen from this analysis that while the computation of T for the 3-D case may appear to be a lengthy computation, each value of T can be computed in a separate parallel process.

Our discussion so far has been limited to determining a line segment that is either rejected or clipped to fit inside a convex set S . Our output (T_L, T_U) can also be used to draw the complement, that is the line segment(s) exterior to a given convex set S .

The possibilities are summarized in Figure 10.

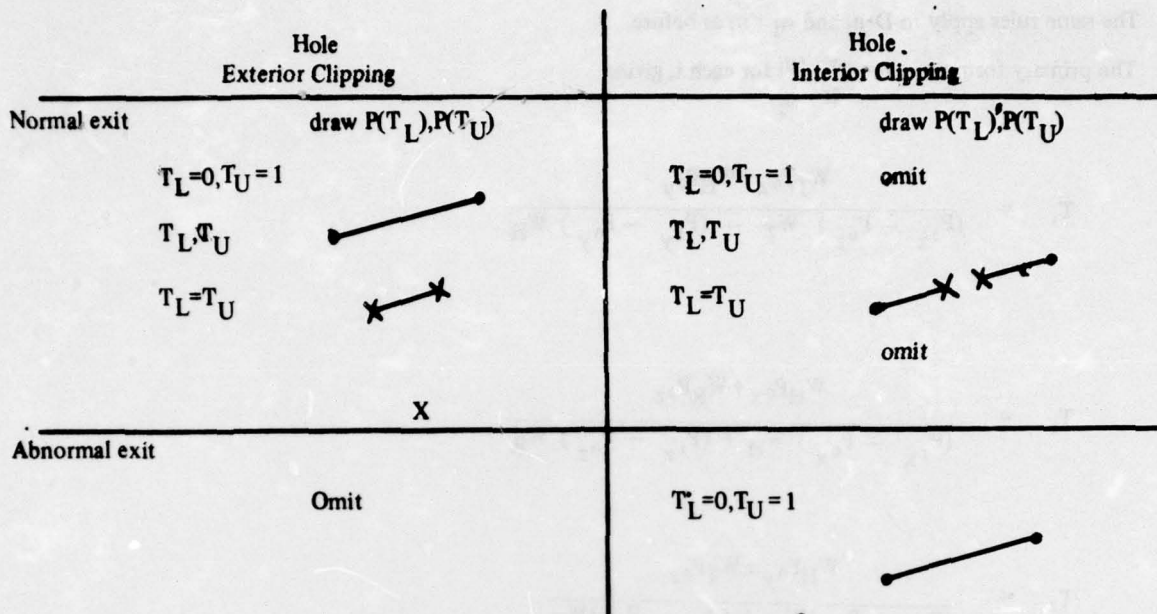


Figure 10. Window-hole clipping summary.

This introduces the possibility that an Algol, PL/1 or Pascal implementation of this algorithm might be implemented to first apply exterior clipping to a 'Hole' contained within the convex window.

If we define a function which performs $P(T)$ calculation, the flow chart (Figure 2) previously drawn can be extended as in Figure 11. The output of SE is scaled appropriately for a display device using integer coordinates.

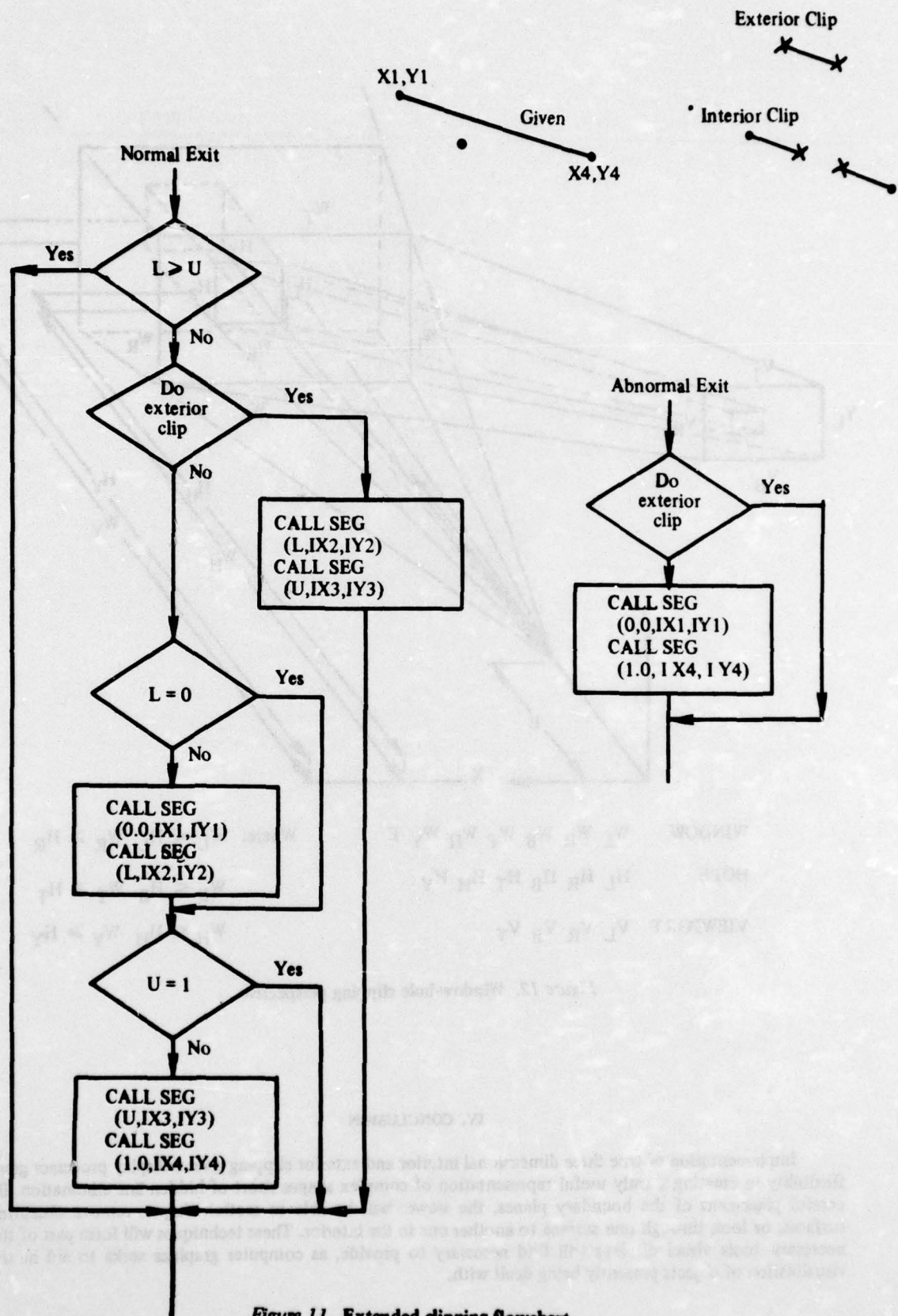
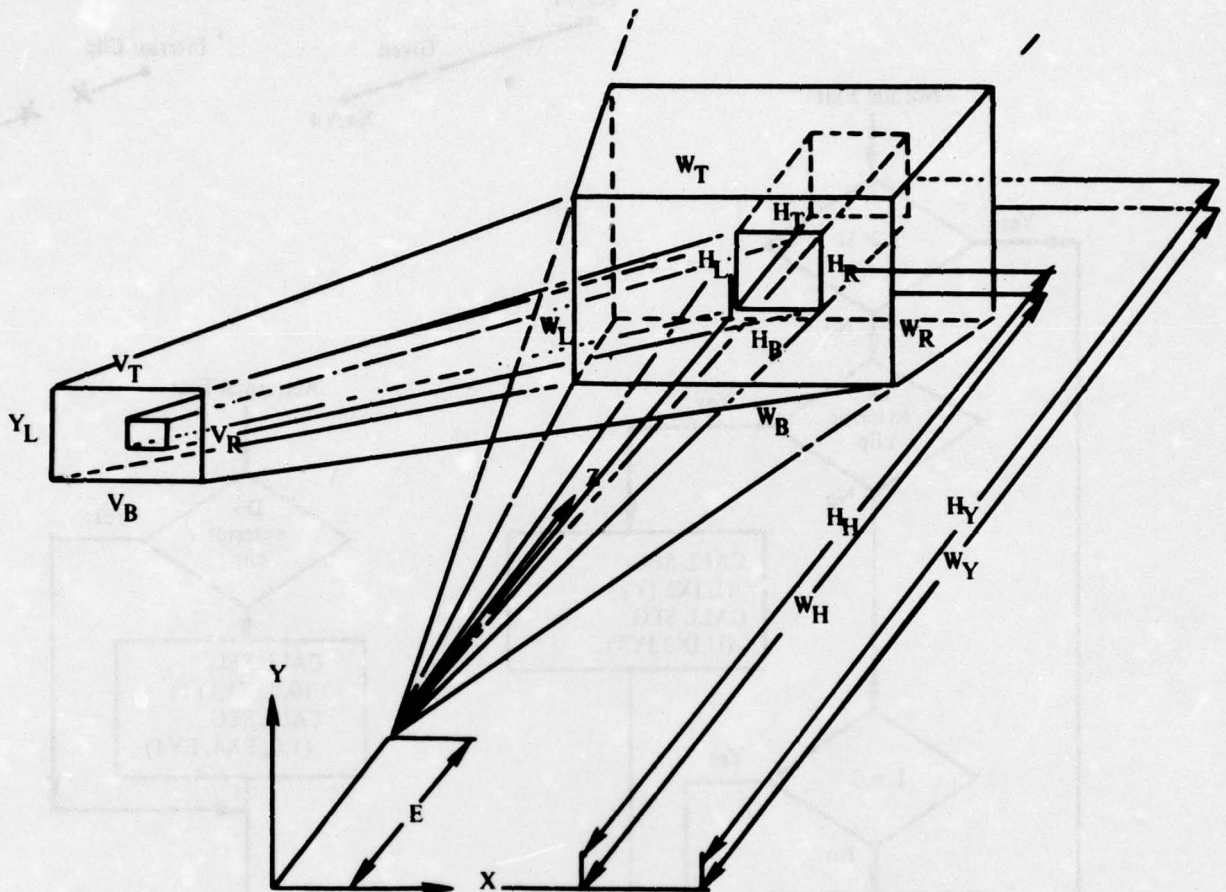


Figure 11. Extended clipping flowchart.



WINDOW $W_L W_R W_B W_T W_H W_Y E$
 HOLE $H_L H_R H_B H_T H_H H_Y$
 VIEWPORT $V_L V_R V_B V_T$

Where: $W_L \leq H_L$ $W_R \geq H_R$
 $W_B \leq H_B$ $W_T \geq H_T$
 $W_H \leq H_H$ $W_Y \geq H_Y$

Figure 12. Window-hole clipping perspective.

IV. CONCLUSION

Implementation of true three dimensional interior and exterior clipping gives a display processor great flexibility in creating a truly useful representation of complex shapes short of hidden line elimination. By careful placement of the boundary planes, the viewer will be able to section shapes, remove obscuring surfaces, or look through one surface to another one in the interior. These techniques will form part of the necessary tools visual displays will find necessary to provide, as computer graphics seeks to aid in the visualization of objects presently being dealt with.